

# Strategic Reasoning for the Resolution of Assignment Problems in Goal- and Actor-Oriented Requirements Engineering

Christophe Chareton

Onera, Toulouse

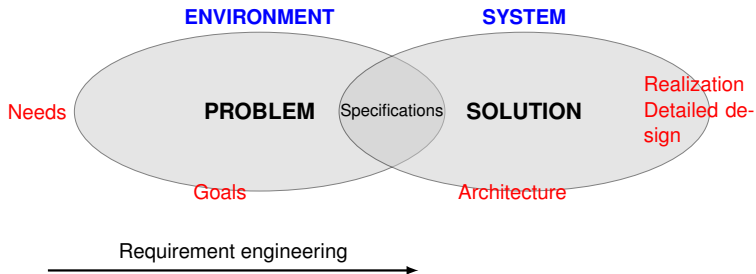
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# Outline

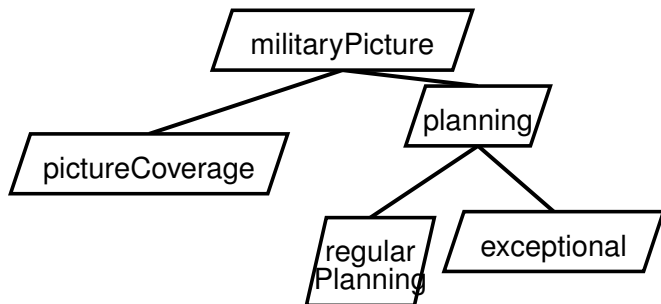
- 1 Language K<sub>H</sub> for Requirements Engineering
- 2 Formalisation
- 3 Assignment problems
- 4 Conclusion : interests of the approach and future works

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# Requirements engineering



- Identify requirements
- Derive specifications from functional requirements
- We concentrate on modeling languages

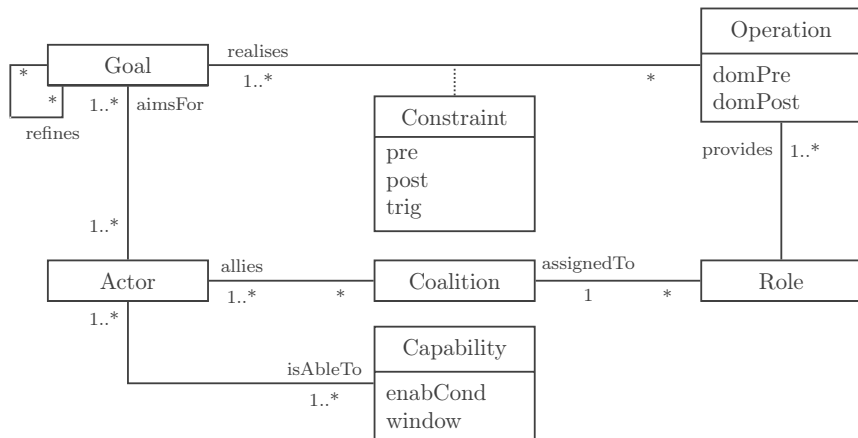


- Formalize the goal's refinement description down to specifications of transitions of the system: KAOS
- Care the realizability of the goals through assignments of roles to actors: TROPOS-I\*

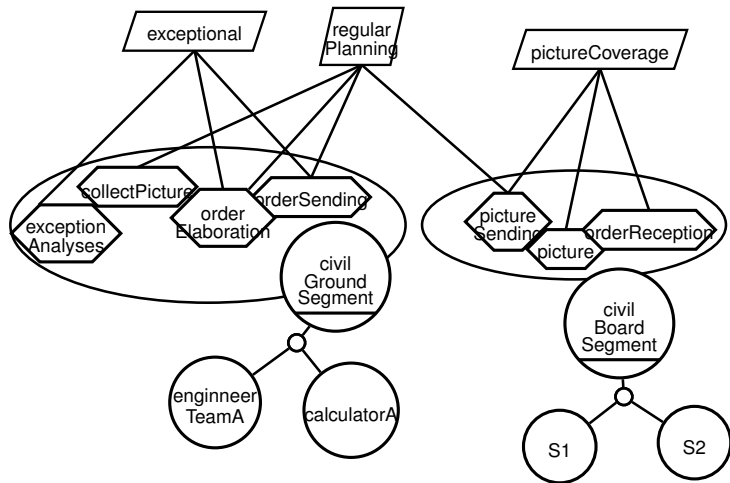
# Motivations for $K_{HI}$ :

- Unify in a single language :
  - The relation pursued goals-operations and its dynamic aspects (KAOS)
  - Intentional agents and the confrontation between available agents (actors) and required agents (roles) (TROPOS-I\*)
- Give a semantic treatment of the capabilities of actors and being able to discuss their ability to play roles: the **assignment problems**

# Metamodel



# Leaf goals, operations, roles and assignment





# Expression of $K_{HI}$ in the logic *USL*

We build USL in three successive steps:

- $\text{Cond}_{K_{HI}}$  : a set of comparisons of values between variables and constants
- The LTL language whose atoms are in  $\text{Cond}_{K_{HI}}$  :  $\text{LTL}_{K_{HI}}$
- Introduction of three operators for strategy treatment:



$$\langle\langle x \rangle\rangle\varphi$$

There is a strategy  $x$  such that  $\varphi(x)$



$$(A \triangleright x)\varphi$$

If agents in  $A$  play along strategy for  $x$  then  $\varphi$



$$(A \not\triangleright x)\varphi$$

If agents in  $A$  do not play anymore along strategy for  $x$  then  $\varphi$

We get a generalization of SL (Fabio Mogavero, Aniello Murano, Giuseppe Perelli, Moshe Y. Vardi)

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# Preliminary definitions

## Definition

- A Non-deterministic Alternating Transition System (*NATS*) is a tuple  $\mathfrak{M} = \langle \Sigma, M, \Pi, \pi, \delta \rangle$  where :
  - $M$  is a set of states, called the domain of the *NATS*,  $\Pi$  is the set of atomic propositions in the language and  $\pi$  is a valuation function, from  $M$  to  $\mathcal{P}(\Pi)$
  - $\delta : \Sigma \times M \rightarrow \mathcal{P}(\mathcal{P}(Q))$  is a transition function mapping a pair  $\langle \text{agent}, \text{state} \rangle$  to a non-empty family of choices of possible next states.
- A *strategy* is a function  $\sigma$  from  $\Sigma \times M$  to  $\mathcal{P}(M)$  such that for all  $(a, s) \in \Sigma \times M, \sigma(a, s) \in \delta(a, s)$
- A *context*  $\kappa$  is a finite word upon  $(\Sigma \times X)^*$ , representing the structure of the active bindings.
- A *memory*  $\mu$  is a partial function from  $X$  to *Strat*, storing the memory instantiations for quantified strategies.

# Satisfaction relation

Let  $\mathcal{M}$  be a NATS, then for all memory  $\mu$ , context  $\kappa$ , state  $s$ :

- $\mathcal{M}, \mu, \kappa, s \models \langle\langle x \rangle\rangle \varphi$  iff there is a strategy  $\sigma \in \mathit{Strat}$  such that  $\mathcal{M}, \mu[x \rightarrow \sigma], \kappa, s \models \varphi$
- $\mathcal{M}, \mu, \kappa, s \models (A, x \triangleright \varphi)$  iff for all  $\lambda$  in  $\mathit{out}(\mu, \kappa[A \rightarrow x])$ ,  $\mathcal{M}, \mu, \kappa[A \rightarrow x]\lambda \models \varphi$
- $\mathcal{M}, \mu, \kappa, s \models (A, x \triangleleft \varphi)$  iff for all  $\lambda$  in  $\mathit{out}(\mu, \kappa[A \leftarrow x])$ ,  $\mathcal{M}, \mu, \kappa[A \leftarrow x]\lambda \models \varphi$

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# Local correctness

Local Correctness is the property of a model such that each of its role can be individually played by its assigned coalition.

## Definition (LC)

Let  $\mathcal{K}$  be an instance of  $\mathcal{KH}_1$ , then the assignment is locally correct, written  $LC(\text{assignedTo})$ , if

$$\mathcal{G}_{\mathcal{K}}, \emptyset, \emptyset, s \models \bigwedge_{r_i \in \text{roles}} \llbracket x_{r_i} \rrbracket (r_i.\text{assignedTo.allies}, x_{r_i}) \llbracket r_i \rrbracket$$

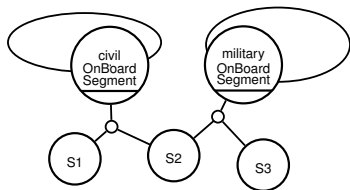
# Global correctness

Global Correctness is the property of a model such that each coalition assigned roles can play them coherently altogether?

## Definition (GC)

Let  $\mathcal{K}$  be an instance of  $\mathcal{K}_{HI}$ , then the assignment is globally correct, written  $GC(assignedTo)$ , if

$$\mathcal{G}_{\mathcal{K}, \emptyset, \emptyset, s} \models \langle\langle x \rangle\rangle \bigwedge_{r \in roles} (r.assignedTo.allies, x) \llbracket r \rrbracket$$



# Positive interaction

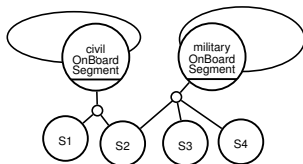
Positive interaction is the relation between roles  $r_1$  and  $r_2$  such that coalition assigned  $r_1$  helps coalition assigned  $r_2$  to play it when itself playing  $r_1$ . Two modalities:

- Possibility

## Definition (PPI)

Let  $\mathcal{K}$  be an instance of  $\text{KHI}$ . Then for all roles  $r_1$  and  $r_2$  in  $\mathcal{K}$ ,  $r_1$  *possibly interacts positively with*  $r_2$ , written  $\text{PPI}(r_1, r_2)$  if

$$\mathcal{G}_{\mathcal{K}}, \emptyset, \emptyset, s \models \langle\langle x \rangle\rangle(r_1.\text{assignedTo.allies}, x)(\llbracket r_1 \rrbracket \wedge \langle\langle y \rangle\rangle(r_2.\text{assignedTo.allies}, y)\llbracket r_2 \rrbracket)$$





# Positive interaction

Positive interaction is the relation between roles  $r_1$  and  $r_2$  such that coalition assigned  $r_1$  helps coalition assigned  $r_2$  to play it when itself playing  $r_1$ . Two modalities:

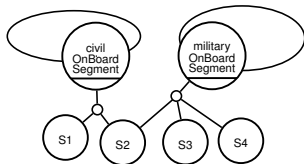
- Necessity

## Definition (NPI)

Let  $\mathcal{K}$  be an instance of  $\text{KHI}$ . Then for all roles  $r_1$  and  $r_2$  of  $\mathcal{K}$ ,  $r_1$  necessarily interacts positively with  $r_2$ , written  $\text{PPI}(r_1, r_2)$  if

$$\mathcal{G}, \emptyset, \emptyset, s \models$$

$$\llbracket x \rrbracket (r_1.\text{assignedTo.allies}, x) (\llbracket r_1 \rrbracket \rightarrow \langle\langle y \rangle\rangle (r_2.\text{assignedTo.allies}, y) \llbracket r_2 \rrbracket)$$



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# Interests of the approach:

Our language gives ...

- ... a semantic for:
  - The dynamics of operations inherited from KAOS
  - A concept of intentional actors : each one pursues its own goals
  - A distinct concept of roles : agents as required entites
  - The expression of several problem of assignments problems
- ... an algorithm to check the relative assignment solutions

## Further works

- Identify lacking actors : in case there is no possible assignment of roles to present actors, non-assignable roles are the actors to introduce by the machine.
- Further characterizations of strategies:
  - Ensure a role  $rl$  assigned to an actor  $a$  does not contradict its pursued goals
  - Compare the efficiency of different strategies in case they do not fully ensure the satisfaction of the roles.

Thank you for your attention

Any question?